

Calculation of high-order virial coefficients for the square-well potential

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(Received 12 April 2016; published 5 July 2016)

Accurate virial coefficients $B_N(\lambda, \varepsilon)$ (where ε is the well depth) for the three-dimensional square-well and square-step potentials are calculated for orders $N = 5-9$ and well widths $\lambda = 1.1-2.0$ using a very fast recursive method. The efficiency of the algorithm is enhanced significantly by exploiting permutation symmetry and by storing integrands for reuse during the calculation. For $N = 9$ the storage requirements become sufficiently large that a parallel algorithm is developed. The methodology is general and is applicable to other discrete potentials. The computed coefficients are precise even near the critical temperature, and thus open up possibilities for analysis of criticality of the system, which is currently not accessible by any other means.

DOI: [10.1103/PhysRevE.94.013301](https://doi.org/10.1103/PhysRevE.94.013301)

I. INTRODUCTION

The pressure P of a fluid at temperature T can be expanded in powers of the number density ρ as

$$\frac{P}{k_B T} = \rho + \sum_{N=2}^{\infty} B_N(T) \rho^N, \quad (1)$$

where k_B is the Boltzmann constant and B_N is the virial coefficient of order N . Despite the fact that this equation of state has a well-established theoretical foundation and is important in the physical sciences, its applications are still rather limited. This is partly because high-order virial coefficients are needed to investigate the convergence properties of the series, and to improve the convergence by, for example, resummation via rational functions or other forms [1–8], but such coefficients are difficult to calculate. Much work has been devoted to the hard-sphere model, for which orders $N \leq 10$ have been calculated [9–11], before a significant improvement in the algorithm was made [12] that gave access to higher orders [12–14]. Recently, the algorithm also made possible calculations beyond $N = 8$ for the Lennard-Jones potential [15].

The square-well (SW) potential is perhaps the simplest model that exhibits behavior that is common to most fluids, including vapor-liquid equilibrium and a critical point. It has been employed to investigate a variety of interesting problems, including interfacial phenomena, surface adsorption, wetting, and capillary condensation [16–18]. The temperature dependence of all virial coefficients of the SW potential can be expressed in an exact closed form (polynomial) and the calculation of the virial coefficients is more tractable than other realistic potentials. The virial coefficients of the SW potential have been studied for more than half a century [19–26]. However, the maximum order attainable through calculation only reached $N = 6$ quite recently [26], which emphasizes the difficulties in obtaining higher-order virial coefficients for model fluids in general, and for the SW potential in particular, relative to the hard-sphere case.

In this paper, we propose an efficient algorithm that enables us to obtain the virial coefficients $B_N(\lambda, f)$, of orders $N = 5$ to 9 for the three-dimensional, pairwise-additive, square-well potential E , which is defined by $E = +\infty$ for an interparticle separation $r \leq \sigma$, $E = -\varepsilon$ for $\sigma < r < \lambda\sigma$, and $E = 0$ for $r \geq \lambda\sigma$. The hard-core diameter is σ . The virial coefficients depend on the reduced well depth $\varepsilon/k_B T$, which is represented here by f , the Mayer function evaluated in the well region [$f = \exp(\varepsilon/k_B T) - 1$], and on the reduced well width λ . The calculations cover the range $1.1 \leq \lambda \leq 2.0$, for which thermodynamic properties such as the critical temperatures and densities are available for comparison. The focus of this paper is primarily on method development and two similar but distinct Monte Carlo methods are discussed.

II. METHODS

The calculated virial coefficients are expressed in reduced form as $B_N^* = B_N/[B_2(\text{HS})]^{N-1}$, where $B_N(\text{HS})$ is the N th virial coefficient for the hard-sphere core, with diameter σ , and $B_2(\text{HS}) = \frac{2\pi}{3}\sigma^3$. As a function of f , the virial coefficients take their hard-sphere values at $f = 0$ (no well): $B_N^*(\lambda, 0) = B_N(\text{HS})/[B_2(\text{HS})]^{N-1} = B_N^*(\text{HS})$. Negative values of f , $-1 \leq f < 0$, correspond to a square-step potential with positive potential energy, with $f = -1$ representing an infinitely high step, for which $B_N^*(\lambda, -1) = \lambda^{3(N-1)} B_N^*(\text{HS})$. Increasingly positive values of f correspond to square-well potentials of greater well depth, and for $f > f_c$ a bulk liquid phase is possible [27], where $f_c = \exp(1/T_c^*) - 1$ and T_c^* is the reduced critical temperature.

Numerical integration (Monte Carlo) is used to calculate the virial coefficients, using the standard integral

$$B_N(\lambda, f) = \frac{1-N}{N!} \int \cdots \int f_B(\lambda, f, \mathbf{r}^N) d\mathbf{r}_2 \cdots d\mathbf{r}_N, \quad (2)$$

and a recursive method of calculation [12] is modified (see Supplemental Material [28]) to evaluate the integrand f_B , which is defined as $f_B(\lambda, f, \mathbf{r}^N) = \sum_G [\prod_{ij \in G} f(r_{ij})]$, where G is a biconnected graph. The integrand, and hence the virial coefficients, are terminating polynomials in f [23,26], which enables the thermodynamic properties of all square-well and

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square-step potentials to be obtained in closed form, for each well width considered, at all temperatures and densities for which the virial series converges. The polynomial expansion of the reduced virial coefficients is given in terms of the coefficients $B_{N,j}^*(\lambda)$ as

$$B_N^*(\lambda, f) = \sum_{j=0}^m B_{N,j}^*(\lambda) f^j. \quad (3)$$

The upper summation limit $m = N(N - 1)/2$ is the number of pairs of particles. From the limiting cases above, it follows that $B_{N,0}^* = B_N^*(\text{HS})$ and

$$\sum_{j=0}^m B_{N,j}^*(-1)^j = \lambda^{3(N-1)} B_N^*(\text{HS}). \quad (4)$$

The recursive method [12] takes as its starting point the quantity $\exp(-E/k_B T)$ for a set of N particles and for all subsets thereof. For the pair-additive square-well potential, this quantity is either zero, if any hard-core overlaps occur in the (sub)set, or $(f + 1)^p$, if p pairs of particles are in the well region and the rest are outside it. The implementation of the recursive method is modified to retain this explicit polynomial dependence on $(f + 1)$ at each stage. The resulting integrand f_B is integrated stochastically to give B_N as a polynomial in $(f + 1)$, which is then converted to the required polynomial in f .

The computer time required for the recursive algorithm scales as $3^N N$ when numerical values are used, but since polynomials of order up to $m = N(N - 1)/2$ are required in this work, the time increases further. It is therefore highly beneficial to invoke the method on a configuration of particles only when needed. One strategy to this end is to store the integrand for reuse during the calculation for each λ and N . The integrand depends only on the graph that describes whether each pair separation is in the hard-core region, the well, or the zero-energy long-range region, and not explicitly on the particle positions within these regions. Given that there are three possibilities per pair, there is a maximum of 3^m possible graphs. For $N = 5$ and 6 ($m = 10$ and 15) the number of graphs is sufficiently small that all the integrands f_B can be stored when first needed, and looked up thereafter.

Two similar but distinct Monte Carlo methods, designated A and B, are implemented to compute SW virial coefficients within this framework, each following respective approaches analogous to previous work for hard spheres [12,14]. Both execute a repeated process of (1) generating a configuration at random by placing each particle in succession relative to a previously positioned particle, starting with the first at the origin; (2) evaluating some simple metrics for the resulting configuration to determine whether to compute f_B for it; (3) if so indicated, computing f_B using the recursive method, and also computing the total probability w for producing the configuration, considering how step (1) could yield it for all $N!$ permutations of the particles; (4) incrementing an average with f_B/w .

In Method A [12], the particles are placed in a chain, such that each new particle is inserted relative to the one just preceding it, with probability that is uniform within the sphere for distances r from 0 to $\lambda\sigma$, and decreasing proportionally

to r^{-12} for larger distances; this probability tail is needed so that all biconnected diagrams have a nonzero chance of being produced. Dynamically generated lookup tables are used to evaluate the integrand.

For $N \geq 7$, the invariance of the integrand to permutation of particle labels is used, giving a storage requirement of order up to $3^m/N!$. To exploit the permutation symmetry, the particle labels should be reordered to produce a canonical labeling; that is, the reordering procedure should produce the same end result regardless of the initial labeling of particles. In Method A, the reordering is done as follows. An adjacency matrix is created with elements A_{ij} equal to 1 if $i = j, 2/3$ if the separation between i and j is in the core region, $1/3$ if it is in the well, and 0 if it is outside the well. The eigenvalues and eigenvectors of the adjacency matrix are calculated, the eigenvector coefficients are squared, and degenerate eigenvectors are added together. The particle with the highest (squared) coefficient in the eigenvector with highest eigenvalue is labeled 1, the one with the second highest coefficient is labeled 2, and so on. Two particles with equal coefficients in the first eigenvector are ranked using their coefficients in the other eigenvectors, in descending order of eigenvalue, and if these are all equal, then the ranking is based on the closest connection to already labeled particles, particle 1 first, using the adjacency matrix.

This reordering is not fully canonical (it does not recognize the permutation equivalence of some graphs), but the number of reordered graphs that are generated is found to be less than $3^m/N!$ in practice. If a graph is not biconnected (where a pair with separation in the core or well region is counted as a connected pair) then the integrand is zero, and is not stored. Many graphs are never produced in practice because they are geometrically impossible (or highly unlikely) in three dimensions. If the storage space is insufficient to accommodate a few rarely generated configurations, this is also acceptable, because the time taken to apply the recursive algorithm to the unstored graphs every time they are generated is not prohibitive. However, for $N = 9$ the storage requirements become sufficiently large that a parallel algorithm is used. Each parallel task is assigned a similar number of biconnected canonical graphs. Each configuration that is generated is checked for biconnectivity, its total probability and canonical ordering are calculated, and these data are sent to the appropriate task for analysis. This balances the storage and processing load between the tasks, and it is found that the computer time is roughly equally divided between the generation of random chains, checking them for biconnectivity, calculating the total probability of generating the chain configurations, performing the canonical reordering, and computing the integrands. More details on this are given in the Results section.

Method B [14] generates configurations via sequential insertion, following tree and ring templates in addition to the chain used by Method A; the insertion probability does not include the r^{-12} tail employed by Method A, because the combination of ring, tree, and chain templates is sufficient to ensure a nonzero probability of generating each biconnected graph without it. The choice of template is made at random, with the tree, ring, or chain weight adjusted in preliminary runs to reduce the variance of the average. Averages of f_B/w are collected in bins, with each bin defined by 5–12 metrics (depending on N ; see Supplemental Material [28])

and uniquely specified for the generated configuration. The estimate of the integral in Eq. (2) depends on the frequency that each bin is visited and the average of f_B/w for all graphs associated with each bin. Ideally, all graphs in a bin will have the same value, such that the variance within the bin is zero. In practice each bin has a nonzero variance (except for $N = 4$, for which the algorithm is equivalent to using a look-up table). The choice of whether to compute f_B/w for a generated configuration is made with probability that depends on the observed variance for its bin, and other factors (with some nonzero probability to compute it regardless of its bin variance), and is adjusted to optimize the calculation [14].

The accuracy of the calculated virial coefficients depends on the quality of the random number generator used, so several different pseudorandom and quasirandom number generators are investigated for producing the random chains. Most of the results presented for Method A are produced using digit-permuted Halton numbers [29], but the Mersenne Twister [30] and random Latin Hypercube [31] methods give consistent results with similar confidence limits; for Method B, the Mersenne Twister is used exclusively.

Uncertainties represent 68% confidence intervals of the corresponding coefficient. For a given N , the coefficients $B_{N,j}^*$ are correlated, so the uncertainty in B_N^* cannot be obtained from them via simple error propagation. Hence, the uncertainties in B_N^* using Method A are computed from averages of the B_N^* at each temperature (f) value reported. Calculations using Method B recorded the covariances needed to propagate the $B_{N,j}^*$ errors to B_N , and these values are included in the Supplemental Material [28].

III. RESULTS AND DISCUSSIONS

Results for the coefficients $B_{N,j}^*$ with well width $1.1 \leq \lambda \leq 2.0$ along with the minimum-uncertainty averages of the two methods are given in the Supplemental Material [28]. A comparison of the uncertainty between the two methods is also provided in this document. For completeness, formulas and results for $N = 2, 3$, and 4 are also included in the Supplemental Material [28].

As outlined above, there are two limiting cases of the SW potential where data for the HS model can be extracted from the calculations: when $f = 0$ [$B_N^*(\text{HS}) = B_N^*(\lambda, 0)$] and $f = -1$ [$B_N^*(\text{HS}) = B_N^*(\lambda, -1)/\lambda^{3(N-1)}$]. These are used as a partial check on the calculated $B_{N,j}^*$. For the fifth virial coefficient with $\lambda = 1.5$, the $B_{5,0}^*$ value is 0.110 253 0(22), and $B_N^*(\text{HS})$ obtained using $f = -1$ is 0.110 251 7(2), which agree well with the literature value 0.110 251 47(6) [9]. The coefficients $B_N^*(\text{HS})$ from the limiting case $f = -1$ always have better precision than those from $f = 0$, because the “large spheres” with diameter $\lambda\sigma$ ($f = -1$) are sampled better than the “small spheres” with diameter σ ($f = 0$). For $N > 5$, there is similar agreement between the HS limiting cases and the HS data from the literature. Thus, these give confidence in the accuracy of all other $B_{N,j}^*$ coefficients reported here.

The temperature dependence of the virial coefficients can be examined by plotting $B_N^*(\lambda, f)$ against f . This is shown in Figs. 1 and 2 for $B_6^*(2, f)$ and $B_9^*(2, f)$ as examples. Similar figures for other $B_N^*(\lambda, f)$ can be generated from the $B_{N,j}^*$ data given in the Supplemental Material [28]. Negative f

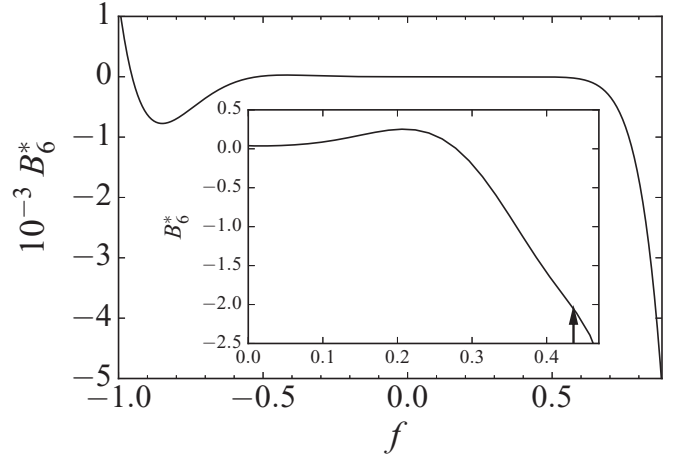


FIG. 1. Sixth virial coefficient of $\lambda = 2.0$. Inset shows an expansion of the supercritical fluid region and the arrow indicates the critical point (f_c). The standard errors are smaller than the width of the lines in the plots.

values correspond to a square-step potential with increasing step height for more negative f , while small positive f values correspond to a SW potential at supercritical temperatures, and large positive f values correspond to low temperatures; $B_N^*(\lambda, f)$ goes to $-\infty$ as f approaches ∞ . For $N = 6$ and $\lambda = 2.0$ (Fig. 1), the standard error in the virial coefficient is very small, which suggests that the calculated $B_6^*(2, f)$ is useful for a wide range of well depths. The standard errors increase from about 1% to 10% for all well widths $\lambda = 1.1-2.0$ near the critical point between order $N = 5$ and 8. For $N = 9$, the uncertainty in $B_9^*(2, f)$ (Fig. 2) is more significant around the critical point, with a percentage error of $\sim 36\%$.

The longest calculation, which is for $B_9^*(1.5, f)$, used a total of 5×10^{12} configurations with Method A and took $\sim 34\,000$ s real time per core when run in parallel on

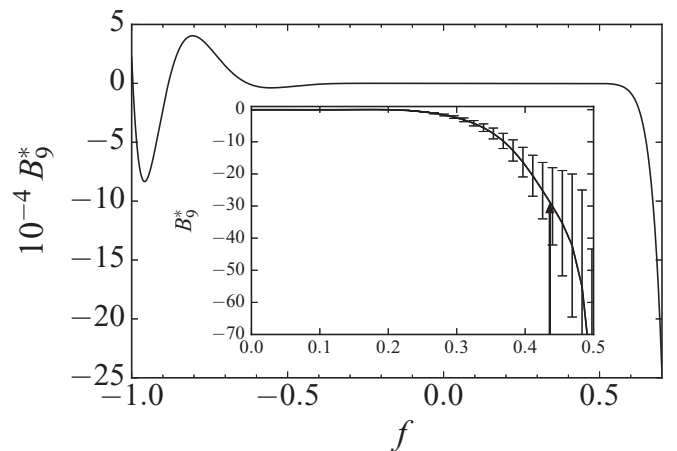


FIG. 2. Ninth virial coefficient of $\lambda = 2.0$. The standard errors are smaller than the width of the line in the main plot. Inset shows an expansion of the supercritical fluid region with error bars denoting the 68% confidence level, and the arrow indicates the critical point (f_c).

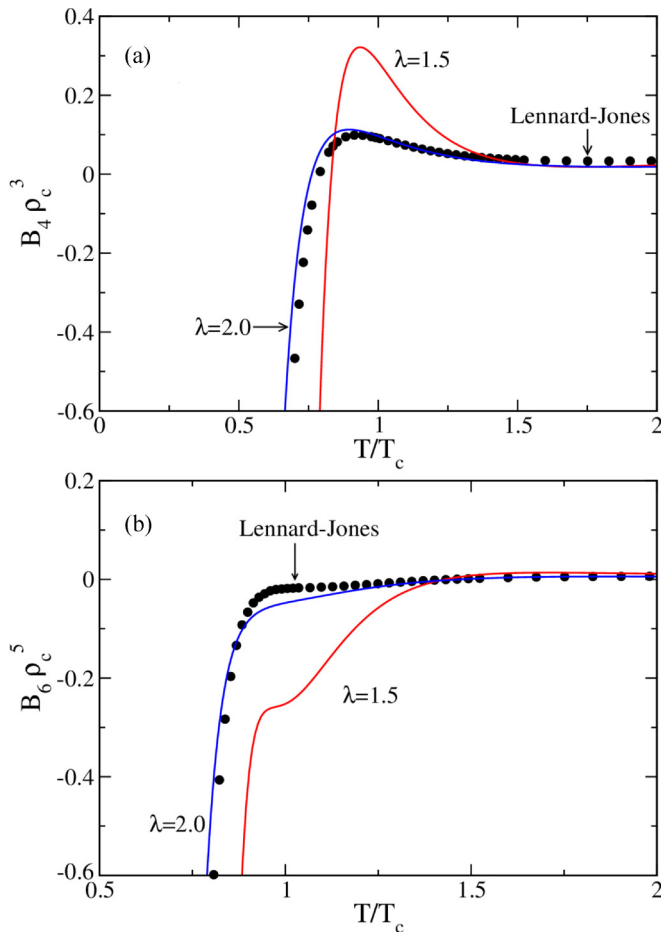


FIG. 3. Contribution of (a) fourth and (b) sixth virial coefficients to the compressibility factor at the critical density, shown for several models as a function of temperature reduced by the critical temperature for each model. Points are for the LJ model [10], and solid lines are results reported here for the SW model with $\lambda = 2.0$ and 1.5 , as indicated. Uncertainties are smaller than symbols and lines.

1200 2.7 GHz CPU cores on a Cray XC30 supercomputer. The number of biconnected graphs found was $\sim 4 \times 10^{11}$. The number of integrands (polynomials) stored was $\sim 3 \times 10^9$, which means that the number of integrands that has to be calculated using the recursive method is reduced by a factor of more than 100 by storing the integrands. The real times required for generating random configurations, checking for biconnectivity, calculating the probability, performing

the canonical reordering, and computing the integrands are $\sim 10\,000$, ~ 1500 , ~ 3000 , $\sim 11\,000$, and ~ 7000 s, respectively, which amounts to $\sim 96\%$ of the total time. The time taken to send and receive data from other processes is therefore insignificant.

Some qualitative differences in behavior of the SW model with respect to λ can be uncovered by comparison of their temperature-dependent virial coefficients when reduced by critical properties. Figure 3 shows the temperature dependence of B_4 and B_6 for $\lambda = 1.5$ and 2.0 , with behavior of the Lennard-Jones (LJ) model provided for reference. The value of the reduced coefficients—which represents their contributions to $P/\rho k_B T$ at the critical density—for $\lambda = 1.5$ is 3.0 and 5.6 times larger (for $N = 4$ and 6 , respectively) than that for $\lambda = 2.0$ when evaluated at their respective critical temperatures. The $\lambda = 2.0$ SW coefficients are much more in line with the behavior of the LJ model, but still the $\lambda = 2.0$ B_6 contribution is 2.6 times greater than that for LJ at their respective critical points. This behavior for $N = 4$ and 6 is illustrative of the other coefficients as well. Accordingly, attempts to evaluate the critical properties for the SW model from these coefficients do not succeed as well as for the LJ model [5]. The use of approximants [5,7] may improve this outcome, and the coefficients reported here should be very useful in formulating such treatments.

IV. CONCLUSION

To conclude, the algorithm developed in this work is capable of obtaining high-order virial coefficients for the SW potential efficiently and accurately. It could also be applied to compute higher-order virial coefficients for other discrete potential models, including the hard-sphere model. Extension to mixtures and systems of different dimensionality is straightforward.

ACKNOWLEDGMENTS

The UK authors are grateful to the University of Nottingham, the MidPlus consortium, and the UK National Supercomputing Computing Service for granting time on their facilities, and the Leverhulme Trust (Grant No. RPG-2614-326) for funding. The US authors thank the University at Buffalo Center for Computational Research for computing resources, and acknowledge funding support from the US National Science Foundation (Grants No. CHE-1027963 and No. CBET-1510017).

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- [28] See Supplemental Material at <http://link.aps.org/supplemental/10.1103/PhysRevE.94.013301> for (a) tabulation of coefficients, (b) method B covariances, (c) graphical comparison of results from two methods, (d) calculation of B_2, B_3 , and B_4 , (e) bin definitions for Method B, and (f) algorithm for calculation of square-well virial coefficients.
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